

Introduction to Numerical Methods in Linear Algebra

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Analysis**
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Outline

- Review last class
- Introduction to numerical methods
- Finite representation of numbers
- Error propagation and round-off error
- Practical considerations in solutions
- Use of pivoting to improve accuracy in solutions of simultaneous linear algebraic equations

Review Eigens

- Eigenvalues and eigenvectors: $\mathbf{Ax} = \lambda\mathbf{x}$
- Computations using $\text{Det}[\mathbf{A} - \lambda\mathbf{I}] = 0$
- Eigenvectors determined by $[\mathbf{A} - \lambda\mathbf{I}]\mathbf{x} = \mathbf{0}$ only to a multiplicative constant
- Define \mathbf{X} as matrix where each column is an eigenvector
- Transformations with a matrix of eigenvectors: $\Lambda = \mathbf{X}^{-1}\mathbf{AX}$

Review Other Matrices

- Orthogonal matrices had orthonormal rows and orthonormal columns
- The inverse of an orthogonal matrix is its transpose
- The eigenvalues of a Hermitian matrix ($\mathbf{A}^* = \mathbf{A}^T$) are real
- An $n \times n$ Hermitian matrix has n linearly independent, orthogonal eigenvalues that can form a unitary matrix

Computer Representations

- Computer is binary machine
 - Numbers represented as series of zeros and one
 - Basic forms are integers and floating point
 - Integers numbers have small range, but exact representation, used for counting
 - Floating point numbers have wide range, but inexact representation
 - Accuracy and range depend on word size

Representing Integers

- Represented as binary number with offset for negative numbers
- Typical computer uses 32 bits (4 bytes) for integer giving range of 0 to $2^{32} - 1$
- Offsets give range from -2^{31} to $2^{31} - 1$
- Adding one to maximum integer gives minimum integer: $(2^{31} - 1) + 1 = -2^{31}$
- Different computers/compiler have different sizes and signed/unsigned

Floating Point Numbers

- Typical sizes are 4 bytes for single precision and 8 for double precision
- Number has sign bit, exponent (characteristic) and mantissa
- IEEE 754 (1985) standard for double
 - 11 bits for exponent
 - 1 sign bit
 - 53 effective bits for mantissa because leading one is not stored

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Machine Epsilon

- Smallest ϵ value such that $1 + \epsilon \neq 1$
- Depends on mantissa bits
- Usually 1.19×10^{-7} for single and 2.22×10^{-16} for double
 - Single almost seven significant figures
 - Double has fifteen-plus significant figures
- Is 10 times $0.1 = 1$? Maybe
 - $0.1_{10} = 0.0001100110011001100110011 \dots_2$
 - $0.1_{10} = 1.100110011001100110011_2 \times 10_2^{-100_2}$

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Rounding Error

- Due to inexact representation
- Store floating point representation \tilde{x} of actual number x with error ϵ_x
- Errors propagate in multiple calculations

$$x = \tilde{x} + \epsilon_x \quad y = \tilde{y} + \epsilon_y$$

$$x \pm y = \tilde{x} + \epsilon_x \pm (\tilde{y} + \epsilon_y) = \tilde{x} \pm \tilde{y} + (\epsilon_x + \epsilon_y)$$

$$\epsilon_{rel} \equiv \frac{(x \pm y) - (\tilde{x} \pm \tilde{y})}{x \pm y} = \frac{\epsilon_x + \epsilon_y}{x \pm y} \approx \frac{\epsilon_x + \epsilon_y}{\tilde{x} \pm \tilde{y}}$$

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Multiplication/Division Error

- Have same result for each although definitions and derivations are different
 - Definitions

$$\epsilon_{rel} = \frac{xy - \tilde{x}\tilde{y}}{xy} \quad \epsilon_{rel} = \frac{\frac{x}{y} - \frac{\tilde{x}}{\tilde{y}}}{\frac{x}{y}}$$
 - Common result $\epsilon_{rel} = \frac{\epsilon_x}{\tilde{x}} + \frac{\epsilon_y}{\tilde{y}}$

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Avoiding Round-off Error

- Use higher precision data types
- Beyond single and double there is extended precision with some compilers
- Systems of equations with large number of equations or nearly singular systems can require double precision
- Algorithms can be designed to reduce round-off error

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Quadratic Real Roots

```

disc = b * b - 4 * a * c;
if ( disc >= 0 )
{ // real solution here
  if ( b > 0 )
    x1 = (-b - sqrt(disc) ) / ( 2 * a );
  else
    x1 = (-b + sqrt(disc) ) / ( 2 * a );

  x2 = c / ( a * x1 );
}
    
```

$$x_1 x_2 = \frac{c}{a}$$

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Gauss Elimination

- Covered previously while discussing solution of systems of equations
- Analytical tool for determining linear dependence or independence
- Basic idea is to manipulate the equations (or data) to make them easier to solve without changing the results
- Systematically create zeros in lower left part of the equations (or data)

Upper Triangular Form

- Convert original set of equations to

$$\begin{bmatrix} \alpha_{11} & \alpha_{12} & \alpha_{13} & \cdots & \cdots & \alpha_{1n-1} & \alpha_{1n} \\ 0 & \alpha_{22} & \alpha_{23} & \cdots & \cdots & \alpha_{2n-1} & \alpha_{2n} \\ 0 & 0 & \alpha_{33} & \cdots & \cdots & \alpha_{3n-1} & \alpha_{3n} \\ \vdots & \vdots & \vdots & \ddots & & & \\ \vdots & \vdots & \vdots & & \ddots & & \\ 0 & 0 & 0 & \cdots & \cdots & \alpha_{n-1n-1} & \alpha_{n-1n} \\ 0 & 0 & 0 & \cdots & \cdots & 0 & \alpha_{nn} \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \\ x_3 \\ \vdots \\ \vdots \\ x_{n-1} \\ x_n \end{bmatrix} = \begin{bmatrix} \beta_1 \\ \beta_2 \\ \beta_3 \\ \vdots \\ \vdots \\ \beta_{n-1} \\ \beta_n \end{bmatrix}$$

Back Substitution

- Upper triangular form on previous slide is easily solved by back substitution
- $x_n = \beta_n / \alpha_{nn}$
- $x_{n-1} = (\beta_{n-1} - \alpha_{n-1n} x_n) / \alpha_{n-1n-1}$, *et cetera*
- General equation for back substitution

$$x_i = \frac{\beta_i - \sum_{j=i+1}^n \alpha_{ij} x_j}{\alpha_{ii}} \quad i = n-1, n-2, \dots, 1$$

Gauss Elimination Algorithm

- Work on **augmented matrix**
- Can handle several **b** vectors at one time

$$[\mathbf{A}, \mathbf{b}] = \begin{bmatrix} a_{11} & a_{12} & a_{13} & \cdots & \cdots & a_{1m} & b_1 \\ a_{21} & a_{22} & a_{23} & \cdots & \cdots & a_{2m} & b_2 \\ a_{31} & a_{32} & a_{33} & \cdots & \cdots & a_{3m} & b_3 \\ \vdots & \vdots & \vdots & \ddots & & \vdots & \vdots \\ \vdots & \vdots & \vdots & & \ddots & \vdots & \vdots \\ a_{n1} & a_{n2} & a_{n3} & \cdots & \cdots & a_{nm} & b_n \end{bmatrix}$$

General Gauss Elimination

- Use each row from row 1 to row n-1 as the "pivot" row
 - Work on each row below the pivot row
 - Multiply pivot row by $a_{row,pivot} / a_{pivot,pivot}$
 - Subtract result from row "row" to make $a_{row,pivot} = 0$
 - Operation requires subtraction for each column of **A** right of pivot column and for **b**
 - Repeat for each row below pivot (except last)
- Use back substitution for **x** values
- Worry about round-off error that depends on the "condition" of the **A** matrix

Problem Condition

- The condition of a problem is defined as the relative change in result divided by a relative change in input
- A large condition number indicates a problem that may give rise to numerical difficulty
- In solution of one equation in one unknown $f(x) = 0$ a small value of df/dx near the root can cause problems

Solving Nonlinear Equations

- Have a system of N simultaneous nonlinear equations written as $\mathbf{f}(\mathbf{x}) = \mathbf{0}$
- Multivariable Newton's method

$$\mathbf{f}(\mathbf{x}^{(n+1)}) = \mathbf{f}(\mathbf{x}^{(n)}) + \mathbf{J}^{(n)}(\mathbf{x}^{(n+1)} - \mathbf{x}^{(n)}) + \dots$$

$$\mathbf{J}^{(n)} = \frac{\partial \mathbf{f}}{\partial \mathbf{x}_m}^{(n)}$$

$$\begin{bmatrix} f_1(\mathbf{x}^{(n+1)}) \\ f_2(\mathbf{x}^{(n+1)}) \\ f_3(\mathbf{x}^{(n+1)}) \\ \vdots \\ f_N(\mathbf{x}^{(n+1)}) \end{bmatrix} = \begin{bmatrix} f_1(\mathbf{x}^{(n)}) \\ f_2(\mathbf{x}^{(n)}) \\ f_3(\mathbf{x}^{(n)}) \\ \vdots \\ f_N(\mathbf{x}^{(n)}) \end{bmatrix} + \begin{bmatrix} J_{11} & J_{12} & J_{13} & \dots & J_{1N} \\ J_{21} & J_{22} & J_{23} & \dots & J_{2N} \\ J_{31} & J_{32} & J_{33} & \dots & J_{3N} \\ \vdots & \vdots & \vdots & \ddots & \vdots \\ J_{N1} & J_{N2} & J_{N3} & \dots & J_{NN} \end{bmatrix} \begin{bmatrix} x_1^{(n+1)} - x_1^{(n)} \\ x_2^{(n+1)} - x_2^{(n)} \\ x_3^{(n+1)} - x_3^{(n)} \\ \vdots \\ x_N^{(n+1)} - x_N^{(n)} \end{bmatrix} = \begin{bmatrix} 0 \\ 0 \\ 0 \\ \vdots \\ 0 \end{bmatrix}$$

Solving Nonlinear Equations II

- Have to solve N simultaneous linear equations at each iteration with $\mathbf{f}(\mathbf{x}^{(n+1)}) = \mathbf{0}$

$$f_k(\mathbf{x}^{(n+1)}) = f_k(\mathbf{x}^{(n)}) + \sum_{m=1}^N J_{km} (x_m^{(n+1)} - x_m^{(n)}) = f_k(\mathbf{x}^{(n)}) + \sum_{m=1}^N \frac{\partial f_k}{\partial x_m}^{(n)} (x_m^{(n+1)} - x_m^{(n)})$$

$$\sum_{m=1}^N J_{km} (x_m^{(n+1)} - x_m^{(n)}) = \sum_{m=1}^N \frac{\partial f_k}{\partial x_m}^{(n)} (x_m^{(n+1)} - x_m^{(n)}) = -f_k(\mathbf{x}^{(n)})$$

- Iterations are affected by physical system through partial derivatives

Solving Nonlinear Equations III

- Thermodynamic property equations are functions of temperature and density
- Want to define state by other properties (e. g., pressure and entropy)
- Newton's method iteration equations

$$\frac{\partial s}{\partial T}(T^{(n+1)} - T^{(n)}) + \frac{\partial s}{\partial v}(v^{(n+1)} - v^{(n)}) = s_0 - s(T^{(n)}, v^{(n)})$$

$$\frac{\partial P}{\partial T}(T^{(n+1)} - T^{(n)}) + \frac{\partial P}{\partial v}(v^{(n+1)} - v^{(n)}) = P_0 - P(T^{(n)}, v^{(n)})$$

- Solve by Cramer's rule

Solving Nonlinear Equations IV

$$(T^{(n+1)} - T^{(n)}) = \frac{\frac{\partial P}{\partial v} [s_0 - s(T^{(n)}, v^{(n)})] - \frac{\partial s}{\partial v} [P_0 - P(T^{(n)}, v^{(n)})]}{\frac{\partial s}{\partial T} \frac{\partial P}{\partial v} - \frac{\partial s}{\partial v} \frac{\partial P}{\partial T}}$$

$$(v^{(n+1)} - v^{(n)}) = \frac{\frac{\partial s}{\partial T} [P_0 - P(T^{(n)}, v^{(n)})] - \frac{\partial P}{\partial T} [s_0 - s(T^{(n)}, v^{(n)})]}{\frac{\partial s}{\partial T} \frac{\partial P}{\partial v} - \frac{\partial s}{\partial v} \frac{\partial P}{\partial T}}$$

- Large partial derivative makes solution ill conditioned

$$\frac{\partial P}{\partial v}$$

Condition of a Matrix I

$$\text{Condition} = \left| \frac{\frac{\text{Change in result}}{\text{result}}}{\frac{\text{Change in input}}{\text{input}}} \right|$$

- Need norms to represent size
- Use matrix norm that is consistent with definition of vector norm

$$\|\mathbf{A}\| \geq \frac{\|\mathbf{Ax}\|}{\|\mathbf{x}\|} \Rightarrow \|\mathbf{A}\| \|\mathbf{x}\| \geq \|\mathbf{Ax}\|$$

Condition of a Matrix II

- Both \mathbf{Ax} and \mathbf{x} are matrices
- Can use any matrix norm to compute $\|\mathbf{A}\|$

- Choosing infinity norm as vector norm gives row sum

$$\|\mathbf{A}\|_{\infty} = \max_i \sum_{j=1}^n |a_{ij}|$$

- Choosing one norm as vector norm gives column sum

$$\|\mathbf{A}\|_1 = \max_j \sum_{i=1}^n |a_{ij}|$$

Condition of a Matrix III

- Correct and incorrect solutions are \mathbf{x} and $\tilde{\mathbf{x}}$
- Computable error residual, $\mathbf{r} = \mathbf{b} - \mathbf{A}\tilde{\mathbf{x}}$

$$\mathbf{r} = \mathbf{b} - \mathbf{A}\tilde{\mathbf{x}} = \mathbf{A}\mathbf{x} - \mathbf{A}\tilde{\mathbf{x}} = \mathbf{A}(\mathbf{x} - \tilde{\mathbf{x}}) = \mathbf{r}$$

$$(\mathbf{x} - \tilde{\mathbf{x}}) = \mathbf{A}^{-1}\mathbf{r} \quad \Rightarrow \quad \|\mathbf{x} - \tilde{\mathbf{x}}\| = \|\mathbf{A}^{-1}\mathbf{r}\| \leq \|\mathbf{A}^{-1}\|\|\mathbf{r}\|$$

$$\|\mathbf{b}\| = \|\mathbf{A}\mathbf{x}\| \leq \|\mathbf{A}\|\|\mathbf{x}\| \quad \Rightarrow \quad \frac{1}{\|\mathbf{x}\|} \leq \frac{\|\mathbf{A}\|}{\|\mathbf{b}\|}$$

$$\frac{\|\mathbf{x} - \tilde{\mathbf{x}}\|}{\|\mathbf{x}\|} \leq \|\mathbf{A}^{-1}\|\|\mathbf{r}\| \frac{1}{\|\mathbf{x}\|} \leq \|\mathbf{A}^{-1}\|\|\mathbf{r}\| \frac{\|\mathbf{A}\|}{\|\mathbf{b}\|}$$

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Condition of a Matrix IV

- Define the relative size of the residual as $\|\mathbf{r}\| / \|\mathbf{b}\|$ (which we can calculate)

$$\frac{\|\mathbf{x} - \tilde{\mathbf{x}}\|}{\|\mathbf{x}\|} \leq \|\mathbf{A}^{-1}\|\|\mathbf{A}\| \frac{\|\mathbf{r}\|}{\|\mathbf{b}\|} = \kappa(\mathbf{A}) \frac{\|\mathbf{r}\|}{\|\mathbf{b}\|}$$

- Condition number $\kappa(\mathbf{A}) \equiv \|\mathbf{A}\| \|\mathbf{A}^{-1}\|$
- Small is < 10 ; large is about 100 or more
- Expect large condition numbers to create problems in solutions

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Condition of a Matrix V

- Ill conditioning comes from near linear dependence in matrix rows

$$\mathbf{A} = \begin{bmatrix} 1.00001 & 0.99999 \\ 1 & 1 \end{bmatrix} \quad \mathbf{A}^{-1} = \begin{bmatrix} 50000 & -49999.5 \\ -50000 & 50000.5 \end{bmatrix}$$

$$\|\mathbf{A}\|_{\infty} = 2 \quad \|\mathbf{A}\|_1 = 2.00001$$

$$\|\mathbf{A}^{-1}\|_{\infty} = 100000.5 \quad \|\mathbf{A}^{-1}\|_1 = 100000$$

Condition number = 200001 confirming ill conditioning apparent from original \mathbf{A}

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Condition of a Matrix VI

- Equations [B-5] and [B-12] in notes show that changes or errors in \mathbf{A} or \mathbf{b} affect solution through the condition number

$$\frac{\|\delta\mathbf{x}\|}{\|\mathbf{x} + \delta\mathbf{x}\|} \approx \frac{\|\delta\mathbf{x}\|}{\|\mathbf{x}\|} \leq \kappa(\mathbf{A}) \frac{\|\delta\mathbf{A}\|}{\|\mathbf{A}\|}$$

$$\frac{\|\delta\mathbf{x}\|}{\|\mathbf{x}\|} \leq \kappa(\mathbf{A}) \frac{\|\delta\mathbf{b}\|}{\|\mathbf{b}\|}$$

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Reducing Round-off Error

- Example of equation order for problem with known solution: $x = 2/3, y = 1/3$

<i>original order</i>	<i>order reversed</i>
$\begin{bmatrix} 0.00003 & 3 \\ 1 & 1 \end{bmatrix} \begin{bmatrix} x \\ y \end{bmatrix} = \begin{bmatrix} 1.0002 \\ 1 \end{bmatrix}$	$\begin{bmatrix} 1 & 1 \\ 0.00003 & 3 \end{bmatrix} \begin{bmatrix} y \\ x \end{bmatrix} = \begin{bmatrix} 1 \\ 1.0002 \end{bmatrix}$

- Gauss elimination gives

$\begin{bmatrix} 0.00003 & 3 \\ 0 & -9999 \end{bmatrix} \begin{bmatrix} x \\ y \end{bmatrix} = \begin{bmatrix} 1.0002 \\ 1 - 1.0002 \end{bmatrix}$	$\begin{bmatrix} 1 & 1 \\ 0 & 2.9997 \end{bmatrix} \begin{bmatrix} y \\ x \end{bmatrix} = \begin{bmatrix} 1 \\ 0.9999 \end{bmatrix}$
--	--

- Look at effect of precision on solutions

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N = Number of Significant Figures

N	Original problem		Equations reversed	
	x	y	x	y
5	.33333	.70000	.33333	.66667
6	.333333	.670000	.333333	.666667
7	.3333333	.6670000	.3333333	.6666667
8	.33333333	.66670000	.33333333	.66666667

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What Happened?

Original	Rows Reversed
$\begin{bmatrix} 0.00003 & 3 \\ 0 & -9999 \end{bmatrix} \begin{bmatrix} x \\ y \end{bmatrix} = \begin{bmatrix} 1.0002 \\ 1 - \frac{1.0002}{0.0003} \end{bmatrix}$	$\begin{bmatrix} 1 & 1 \\ 0 & 2.9997 \end{bmatrix} \begin{bmatrix} y \\ x \end{bmatrix} = \begin{bmatrix} 1 \\ 0.9999 \end{bmatrix}$

- Problem is in the $1 - 1.0002/0.0003$ term
- Division inaccuracy swamps subtraction
- Try to have large elements on pivot row to avoid such divisions
- Pivoting strategy: Use row with maximum (scaled) element as pivot row

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More on Pivoting

- Can exchange rows or columns to get maximum element on pivot
- Exchanging rows only is easier and often effective
 - Commonly used in earlier programs
 - Exchanging columns changes the identity of the unknowns requiring more computer work to keep track of changes
- Newer software now exchanges both

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Can get help from menu or typing help <item> at command line

Command Window to enter commands and get results

List of current variables and their values

Use up-arrow to get previous commands from command history (can edit and execute again)

Double-click command history here to execute again without edits

Windows for files you write in MATLAB

ans =
Result
10.0179

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Entering Arrays

– Enter a row vector by enclosing data in [] separated by a space

```
>> row = [12 -3 5 7 0]
```

– Enter a column vector by enclosing data, separated by a semicolon (;) in []

```
>> col = [-3; 6; 0; 4]
```

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Entering a Matrix

– Enter matrix data row by row

```
>> A = [1 2; 3 4]
```

– Put spaces between data in the same row

```
A =
```

```
1 2
```

```
3 4
```

– Put a semicolon to start data on next row

```
>> B = [1 2 3; 4 5 6; 7 8 9]
```

- MATLAB uses the ... as a continuation signal
- After the ... hit Enter and continue input of same command on a new line

```
B =
```

```
1 2 3
```

```
4 5 6
```

```
7 8 9
```

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Entering a Matrix II

– Pressing enter after each row of data can be used to enter a matrix

```
B = (Result)
```

```
1 2 3
```

```
4 5 6
```

```
7 8 9
```

– Using semicolons, all data can be placed on one row (see below)

```
>> B = [1 2 3; 4 5 6; 7 8 9]
```

– Continuation is only needed to start new line in the middle of data entry

```
>> B = [1 2 3; 4 5 6; 7 8 9]
```

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MATLAB Linear Solver

- For an $n \times n$ matrix \mathbf{A} , and a $n \times m$ right-hand side matrix \mathbf{b} MATLAB produces a $n \times m$ solution $\mathbf{x} = \mathbf{A}^{-1}\mathbf{b}$ by either of the following commands
 - $\mathbf{x} = \mathbf{A}/\mathbf{b}$
 - $\mathbf{x} = \text{mldivide}(\mathbf{A}, \mathbf{b})$
- Each column of the \mathbf{x} result contains the solution of $\mathbf{Ax} = \mathbf{b}$ for the corresponding column of the \mathbf{b} array

MATLAB Eigenvalues/vectors

- Use function eig for both eigenvalues and eigenvectors
- Basic command is $[\mathbf{V}, \mathbf{D}] = \text{eig}(\mathbf{A})$
 - \mathbf{A} is a square matrix
 - \mathbf{D} is a diagonal matrix giving the eigenvalues of \mathbf{A} on the diagonal (\mathbf{A} matrix)
 - \mathbf{V} is a square matrix whose columns are the eigenvectors of \mathbf{A} (\mathbf{X} matrix)

Excel Data to MATLAB

```
>> A = xlsread('Gaussian.xlsx', 'Data',  
'B6:CW105'); File Name Worksheet  
>> b = xlsread('Gaussian.xlsx', 'Data',  
'CX6:DI105'); Data Range  
>> xExact = xlsread('Gaussian.xlsx',  
'Answers', 'B3:M102');  
>> x = A\b  
>> RMS = sqrt(mean(x - xExact).^2)
```